

# Semantics

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# Review

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# Urgent!!!

## Klausur

\* *NEXT WEEK!!!!* \*

- not in this room!!!

## Location

HSG HS 8

## Semantic Interpretation Rules

$$\left[ \begin{array}{c} \bullet \\ | \\ \alpha \end{array} \right] = \llbracket \alpha \rrbracket$$

$$\left[ \begin{array}{c} \bullet \\ / \quad \backslash \\ \alpha \quad \beta \end{array} \right] = \llbracket \alpha \rrbracket \otimes \llbracket \beta \rrbracket$$

$$\left[ \begin{array}{c} \bullet \\ / \quad \backslash \\ \alpha_i \quad \beta \end{array} \right] = \llbracket \alpha \rrbracket \otimes (\lambda x_i. \llbracket \beta \rrbracket)$$

$$\llbracket t_j \rrbracket = x_j$$

## Scope

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[one [apple [in every basket]]] is rotten

in eet

basket,apple et

one,every (et)(et)t

be (et)et

rotten et

**compare with:**

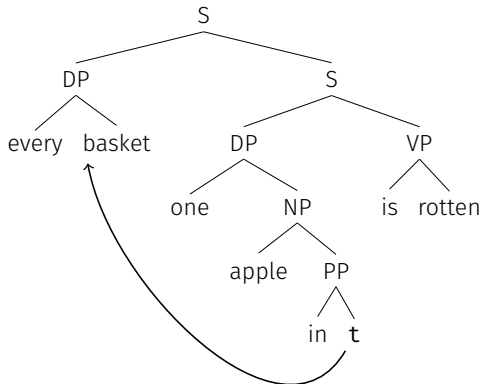
*one apple which was in every basket was rotten*

## Interpreting this sentence

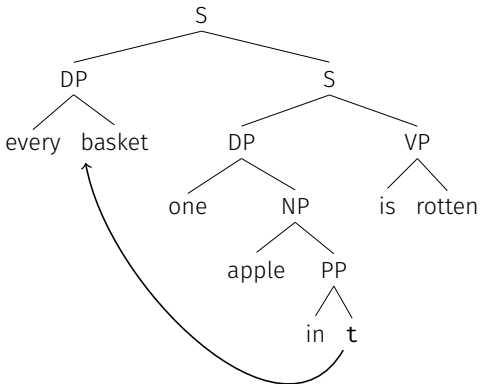
one apple in every basket is rotten

$\text{every}(\text{basket})(\lambda x.\text{one}(\text{apple} \wedge \text{in}(x))(\text{rotten}))$

looks like movement of *every basket* to a higher position



But...



...isn't this an island violation?

\*which basket is one apple in rotten?



## Ambiguous?

### Usually implausible with a single sentence

1. one apple in every basket is rotten
2. no apple in a basket is rotten

### [D [N [P DP]]] VP

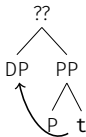
1. for *every basket*, there is *one apple* in *that basket*, such that *that apple* is rotten  
 $DP(\lambda x.D(N \wedge P(x))(VP))$
2. for *no apple* which is in a *basket*, is *that apple* rotten  
 $D(N \wedge \lambda y.DP(\lambda x.P(x)(y)))(VP)$

## How could we build these?

$D(N \wedge \lambda y.DP(\lambda x.P(x)(y)))(VP)$

Focus on  $\lambda y.DP(\lambda x.P(x)(y))$

- ignoring the  $\lambda y$  for the moment
- $DP(\lambda x.P(x)(y))$  looks like it would come from the following tree:



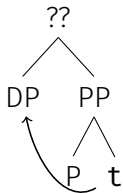
## Putting the $\lambda$ back

doesn't work!

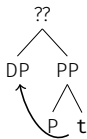
- $P : eet$
- $\lambda x.P(x) : eet$
- $DP : (et)t$

But of course:

- $\lambda x.P(x)(y) : et$



## Putting the $\lambda$ back (II)



$\lambda y. DP(\lambda x. P(x)(y))$

$P = \lambda z. \lambda y. P(z)(y)$

### a movement construction

- have  $[[DP]]$  on one hand,

- and  $\lambda x. \left[ \left[ \begin{array}{c} PP \\ / \quad \backslash \\ P \quad t \end{array} \right] \right] = \lambda x. \lambda y. P(x)(y)$  on the other

how to put them together?

## Putting the $\lambda$ back (III)

$$\text{DP} \oplus \lambda x. \lambda y. P(x)(y) = \lambda y. \text{DP}(\lambda x. P(x)(y))$$

### Adding a new mode of composition

$$f \oplus g = f(g)$$

$$f \oplus g = \lambda y. f(\lambda x. g(x)(y))$$

### Types

- $g : abc$
- $f : (ac)d$
- $\lambda y. f(\lambda x. g(x)(y)) : bd$

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# Bind in the continuation monad

## Some notation

write  $f \gg= g$  for  $\lambda y.f(\lambda x.g(x)(y))$

$GQ \gg= R$

apply a  $GQ$  to the first argument of a binary relation

- prepositions
- transitive verbs (!)

## Putting everything together

What structure will give us the following term?

$D(N \wedge DP \gg= \lambda x.P(x))(VP)$

Can we avoid moving out of the DP?

$DP(\lambda x.D(N \wedge P(x))(VP))$