

Replicative Processes: Morphology

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Non-Local Metathesis in Frampton's (2009) Approach

Background:

Frampton's system allows for local metathesis, given adjacent sequences each of which are enclosed by <, > angle brackets.

Question:

Does Frampton's system permit non-local metathesis where there is intervening material?

Hypothesis:

Yes, it does, given that duplication juncture insertion and truncation juncture insertion can be recursive (see p. 97 on double reduplication in Mokilese).

(1) *Transcription rules* (left)

- a. Init:
 $\emptyset \rightarrow \langle \rangle / \underline{\quad} \times$
- b. Copy Shift:
 $\Delta \dots \times \rightarrow \times \Delta \dots \times \rangle$
- c. Cleanup:
 $[\langle \times^* \rangle] \rightarrow \times^*$
- d. Copy Delete:
 $\Delta \dots \rangle \langle \times \rightarrow \times \Delta \dots \rangle \langle$
- e. $\langle \rangle$ -Cleanup:
 $\langle \rangle \rightarrow \emptyset / \underline{\quad}$

(2) *Local metathesis*

- a. $[\langle a \rangle \langle b \rangle]$ (Copy Shift is not applicable, but Copy Delete is)
- b. $b [\langle a \rangle \langle \rangle]$ ($\langle \rangle$ -Cleanup)

- c. $b [\langle a \rangle]$ (Cleanup)
- d. $b a$

Intended result with non-local metathesis:

$a b c d \rightarrow c b a d$

(3) *Non-local metathesis via transcription juncture recursion*

- a. $[\langle [\langle a \rangle \langle b \rangle] \rangle \langle c \rangle] d$ (Copy Delete)
- b. $[\langle b [\langle a \rangle \langle \rangle] \rangle \langle c \rangle] d$ ($\langle \rangle$ -Cleanup)
- c. $[\langle b [\langle a \rangle] \rangle \langle c \rangle] d$ (Cleanup)
- d. $[\langle b a \rangle \langle c \rangle] d$ (Copy Delete)
- e. $c [\langle b a \rangle \langle \rangle] d$ ($\langle \rangle$ -Cleanup)
- f. $c [\langle b a \rangle] d$ (Cleanup)
- g. $c b a d$

Note:

It is not quite obvious to me how the rules for duplication and truncation juncture insertion in the reduplicant are best formulated. An obvious way to proceed would be to do this via prosodic adjustment (chapter 6), in terms of goal-driven rules (GDR) for the derived conjunct.

An application: Schüttelreime (~ "spoonerisms")

(4) *Bierstube – Stierbube*

b	i: ɤ	f t	u: b ə
a	b	c	d
c	b	a	d
f t	i: ɤ	b	u: b ə

(5) *Ableitung*

- a. $b i: \tilde{\gamma} f t u: b \tilde{\epsilon}$ (juncture insertion)
- b. $[\langle [\langle b \rangle \langle i: \tilde{\gamma} \rangle] \rangle \langle f t \rangle] u: b \tilde{\epsilon}$ (Copy Delete)
- c. $[\langle i: \tilde{\gamma} [\langle b \rangle \langle \tilde{\gamma} \rangle] \rangle \langle f t \rangle] u: b \tilde{\epsilon}$ (Copy Delete)
- d. $[\langle i: \tilde{\gamma} [\langle b \rangle \langle \rangle] \rangle \langle f t \rangle] u: b \tilde{\epsilon}$ ($\langle \rangle$ -Cleanup)
- e. $[\langle i: \tilde{\gamma} [\langle b \rangle] \rangle \langle f t \rangle] u: b \tilde{\epsilon}$ (Cleanup)
- f. $[\langle i: \tilde{\gamma} b \rangle \langle f t \rangle] u: b \tilde{\epsilon}$ (Copy Delete)
- g. $f [\langle i: \tilde{\gamma} b \rangle \langle t \rangle] u: b \tilde{\epsilon}$ (Copy Delete)
- h. $f t [\langle i: \tilde{\gamma} b \rangle \langle \rangle] u: b \tilde{\epsilon}$ ($\langle \rangle$ -Cleanup)
- i. $f t [\langle i: \tilde{\gamma} b \rangle] u: b \tilde{\epsilon}$ (Cleanup)
- j. $f t i: \tilde{\gamma} b u: b \tilde{\epsilon}$

Frampton, John (2009): *Distributed Reduplication*. MIT Press, Cambridge, Mass.